Concurrent Programming: Languages and Techniques

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MIEI - Integrated Masters in Comp. Science and Informatics Specialization Block

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Part I Administrivia



3 Modules:

- 1. Message-passing Concurrency (Go)
- 2. Message-passing Conc. + Distribution (Erlang)
- 3. Memory safe Shared-memory Concurrency (Rust)



Grading:

- 3 Mini-Projects (5% each, 2 best)
- 3 Projects (**20%** each)
- 3 Mini-tests (**10%** each, no min. grade)

Administrivia

Final grade is 70% projects (and mini) + 30% tests



- Projects and Mini-Projects:
- Groups of 2 students (3 only if approved!)
- Mini-projects are 1 week-long small scale
- Projects are 2 week-long medium/small scale



Mini-tests (~1h30 duration):

- After each module, covers material from the module.
- Multiple choice (12 points), focusing more on the concepts / theoretical aspects.
- Open answer (8 points)
- Includes (open answer) questions about the project!
- Sample test will be made available in due time.



- Typical module structure:
- Week 1: Intro to language & prog. paradigm
- Week 2: Mini-Project handout (due following week)
- Week 3: Project handout (due in 2 weeks)
- Week 4: Final lecture and project support



Important dates (tentative):

- Week of Oct-3: Mini-Project 1 Deadline
- Week of Oct-17: Project 1 Deadline
- Week of Oct-24: Mini-Test 1 ----
- Week of Oct-31: Mini-Project 2 Deadline
- Week of Nov-14: Project 2 Deadline
- Week of Nov-21: Mini-Test 2
- Week of Nov-28: Mini-Project 3 Deadline
- Week of Dec-12: Project 3 Deadline
- Week of Dec-19: Mini-Test 3

Part II Language-based Problem Solving

Programming is Hard!

- Programming is about convincing machines to do what we want them to do.
- Systematic and rigorous ("No more, no less").
- The way in which we communicate with machines shapes what we can (or can't) do.
- ...and how easy it is achieve it.

Programming is Hard!

- Programming languages and their features shape what our programs can and cannot do.
- "If the only tool you have is a hammer, you tend to see every problem as a nail."
- Lets think a bit about what kind of "hammers" we have...

What about PLs?

- Two axes:
 - What kind of runtime errors are allowed?
 - language features?

What idioms and paradigms are facilitated by

- Different languages allow programs to go wrong in different ways.
- We distinguish between runtime and compile-time errors:
 - Anything goes (e.g. Python)
 - Memory violation errors (e.g. C/C++)
 - Null pointer dereferencing (e.g. Java)
 - Memory safety (e.g. Rust)



- Beyond errors, languages also provide abstractions that can be better suited for certain problems.
- Languages and their implementations have a very wide range of focus.
- Different problems are better solved using specific (language) features.

Writing a compiler / interpreter / static analyzer?

```
data Type
    = Base
    Arrow Type Type
    deriving (Eq, Ord, Read, Show)
data Term
    = Const
    | Var Int -- deBruijn indexing; the nearest enclosing lambda binds Var 0
    | Lam Type Term
    App Term Term
    deriving (Eq, Ord, Read, Show)
check :: [Type] -> Term -> Maybe Type
check env Const = return Base
check env (Var v) = atMay env v
check env (Lam ty tm) = Arrow ty <$> check (ty:env) tm
check env (App tm tm') = do
    Arrow i o <- check env tm
   i' <- check env tm'
   guard (i == i')
    return o
eval :: Term -> Term
eval (App tm tm') = case eval tm of
Lam _ body -> eval (subst 0 tm' body)
eval v = v
subst :: Int -> Term -> Term -> Term
subst n tm Const = Const
subst n tm (Var m) = case compare m n of
   LT -> Var m
```

VS

Ľ	Expr.java
Ľ	ExprBetaReducer.java
Ľ	ExprBuilder.java
Ľ	ExprEtaReducer.java
Ľ	ExprNone.java
Ľ	ExprParser.java
Ľ	ExprParserUntyped.java
Ľ	ExprPredicate.java
Ľ	ExprPrinter.java
Ľ	ExprRichBuilder.java
Ľ	ExprToDeBruijn.java
Ľ	ExprToFreeNames.java
Ľ	ExprToType.java
L P	FreehNemeSupplierieue

Web development?

```
html>
<head>
   <div>
          <div>
              <form method="post" action="#" id="formvalue" onkeyup="</pre>
              drawChart()" />
              </form>
          </div>
  </div>
  <script type="text/javascript" src="https://www.google.com/jsapi"><//
</pre>
   script>
  <script type="text/javascript">
  var bid = 43;
  var ask = 21;
    google.load("visualization", "1", {packages:["corechart"]});
    google.setOnLoadCallback(drawChart);
    function drawChart() {
      var data = google.visualization.arrayToDataTable([
        ['Price', 'Quantity'],
         ['Value #1', bid],
         ['Value #2', ask],
       ]);
```

.....

type filterFn = (item: string) => bool; // The higher-order-function takes an array and a function as arguments function filterItems(arr: string[], fn: filterFn): string[] { const newArray: string[] = []; arr.forEach(item => { if(fn(item)){ newArray.push(item); }); return newArray; function checkNameLength(name: string) { return name.length >= 10; const doctorList = ["DoctorOne", "DoctorTwo", "DoctorThree", "DoctorFour"]; // We are passing the array and a function as arguments to filterItems method. const output = filterItems(doctorList, checkNameLength); console.log(output); // ["DoctorThree", "DoctorFour"]

Javascript







Elm



Device driver?

```
#include <linux/module.h>
#include <linux/string.h>
#include <linux/fs.h>
#include <asm/uaccess.h>
// module attributes
MODULE_LICENSE("GPL"); // this avoids kernel taint warning
MODULE DESCRIPTION("Device Driver Demo");
MODULE_AUTHOR("Appu Sajeev");
static char msg[100]={0};
static short readPos=0;
static int times = 0;
// protoypes,else the structure initialization that follows fail
static int dev_open(struct inode *, struct file *);
static int dev_rls(struct inode *, struct file *);
static ssize t dev read(struct file *, char *, size t, loff t *);
static ssize_t dev_write(struct file *, const char *, size_t, loff_t *);
// structure containing callbacks
static struct file_operations fops =
        .read = dev_read, // address of dev_read
        .open = dev_open, // address of dev_open
        .write = dev_write, // address of dev_write
        .release = dev_rls, // address of dev_rls
};
// called when module is loaded, similar to main()
int init module(void)
        int t = register_chrdev(89, "myDev", &fops); //register driver with major:89
        if (t<0) printk(KERN_ALERT "Device registration failed..\n");</pre>
        else printk(KERN_ALERT "Device registered...\n");
        return t;
 // called when module is unloaded, similar to destructor in OOP
void cleanup module(void)
```

C/C++

```
/// Communicating to the sensor, updates the latest wrench.
/// # Returns
/// `Ok(wrench)` if succeeds, `Err(reason)` if failed.
pub fn update(&mut self) -> Result<Wrench, Error> {
   let res = receive_message(&mut self.port);
   // Regardless of success or failure of receive_message(), request the next single data.
   // If we do not so, after updating failed once, updating will fail everytime due to no reception from the sensor.
   self.request_next_wrench()?;
   let res = match res {
       0k(res) => res,
       Err(e) => {
           return Err(e);
   };
   let (fx, fy, fz, mx, my, mz) = (0..6)
        .map(|i| 4 + i * 2)
        .filter_map(|start| res.get(start..start + 2))
        .map(|res| i16::from_le_bytes(res.try_into().unwrap()))
        .map(|digital| digital as f64)
        .next_tuple()
        .ok_or(Error::ParseData)?;
   let rated_binary = self.product.rated_binary();
   let force = Triplet::new(fx, fy, fz)
        .map_entrywise(self.rated_wrench.force, |left, right| {
           left / rated_binary * right
       });
   let torque = Triplet::new(mx, my, mz)
        .map_entrywise(self.rated_wrench.torque, |left, right| {
            left / rated_binary * right
       });
   self.last_raw_wrench = Wrench::new(force, torque);
   Ok(self.last_wrench())
```

Rust

Languages as Tools Operating System?

```
/// An ELF executable
pub struct Elf<'a> {
    pub data: &'a [u8],
    header: &'a header::Header
}
impl<'a> Elf<'a> {
    /// Create a ELF executable from data
    pub fn from(data: &'a [u8]) -> Result<Elf<'a>, String> {
       if data.len() < header::SIZEOF_EHDR {</pre>
            Err(format!("Elf: Not enough data: {} < {}", data.len(), header::SIZEOF_EHDR))</pre>
       } else if &data[..header::SELFMAG] != header::ELFMAG {
            Err(format!("Elf: Invalid magic: {:?} != {:?}", &data[..header::SELFMAG], header::ELFMAG))
       } else if data.get(header::EI_CLASS) != Some(&header::ELFCLASS) {
            Err(format!("Elf: Invalid architecture: {:?} != {:?}", data.get(header::EI_CLASS), header::ELFCLASS))
       } else {
            Ok(Elf {
                data,
                header: unsafe { &*(data.as_ptr() as usize as *const header::Header) }
           })
       }
    }
    pub fn sections(&'a self) -> ElfSections<'a> {
        ElfSections {
            data: self.data,
            header: self.header,
            i: 0
       }
    }
    pub fn segments(&'a self) -> ElfSegments<'a> {
        ElfSegments {
            data: self.data,
            header: self.header,
            i: 0
    }
                                             Rust
```

```
module StateTick = State.Make(Time)
(* Individual received TCP segment
   TODO: this will change when IP fragments work *)
type segment = { header: Tcp_packet.t; payload: Cstruct.t }
let pp_segment fmt {header; payload} =
 Format.fprintf fmt
    "RX seg seq=%a acknum=%a ack=%b rst=%b syn=%b fin=%b win=%d len=%d"
    Sequence.pp header.sequence Sequence.pp header.ack_number
   header.ack header.rst header.syn header.fin
    header.window (Cstruct.length payload)
let len seg =
 Sequence.of_int ((Cstruct.length seg.payload) +
 (if seg.header.fin then 1 else 0) +
 (if seg.header.syn then 1 else 0))
(* Set of segments, ordered by sequence number *)
module S = Set.Make(struct
   type t = segment
   let compare a b = (Sequence.compare a.header.sequence b.header.sequence)
 end)
type t = {
 mutable segs: S.t;
  rx_data: (Cstruct.t list option * Sequence.t option) Lwt_mvar.t; (* User receive channel *)
 ack: ACK.t;
  tx_ack: (Sequence.t * int) Lwt_mvar.t; (* Acks of our transmitted segs *)
  wnd: Window.t;
  state: State.t;
```





- Different lang. impl. provide different trade-offs:
 - Efficiency
 - Correctness
 - Rapid prototyping
 - Ease of refactoring
 - Interoperability

will often depend on the tool / language.

<pre>yuvraj@DeathNote:~/hea Search Results:</pre>
William: SBM GSOH like Josh:SJM likes sports, Segmentation fault (co
yuvraj@DeathNote:~/hea

Correctness is mandatory, but how hard it is to get it

d-first-c-exercises\$./find sports, TV, dining novie theatre re dumped) d-first-c-exercises\$

Part II Concurrency



- In a single processor, programs executed "simultaneously"
- programs (processes)
- Processes compete for processor access
- processors

- Scheduler distributes "processor time" to running

- This is true **regardless** of the number of physical

Parallelism vs Concurrency

Parallelism: Programming as the simultaneous execution of (possibly related) computations.

Concurrency: Programming as the composition of independently executing processes.













picture from https://talks.golang.org/

Parallelism vs Concurrency

Parallelism: Programming as the simultaneous execution of (possibly related) computations.

independently executing processes.







- **Concurrency:** Programming as the composition of











Concurrency vs Parallelism

- concurrency!
- Programs can be concurrent and have 0 parallelism.

- Concurrency is not parallelism, but parallelism is enabled by

- Well-written concurrency may run better on a multiprocessor.

Concurrency and Independence

- Concurrency is a way to **structure** work into independent pieces ...

... but then you have to coordinate those pieces

- "Independent" here refers to a way of thinking about problems, and structuring their solutions.

- Concurrent processes may indeed interfere/interact

Andrew Gerrand (Golang)

Type of Concurrency

Shared memory concurrency:

- Processes coordinate by reading and writing to memory locations that are shared.
- Concurrent memory accesses managed by locks.

Message-passing concurrency:

- Processes coordinate by sending and receiving messages along **channels**.
- More abstract / higher-level than shared memory concurrency
- Channel operations need not be managed (processes may concurrently perform channel ops **safely**).

Challenges

- problem *state explosion problem*.
- have 6 possible executions / interleavings.

- Reasoning about possible executions is a combinatoric

- With $P = P_1$; P_2 and $Q = Q_1$; Q_2 , P and Q in parallel

- Errors can be quite difficult to diagnose / reproduce.

Safe Concurrency

- What do we mean by safe? What can go wrong?
- Uncontrolled concurrent accesses to data (Data Races)
- Being stuck forever acquiring a lock (Deadlock)
- Being stuck forever trying to read / write to a channel (Deadlock)
- Repeating the same interaction without doing any useful work (Livelock)

Safe Concurrency

- Correctness properties:
- consistency)
- fully "stuck")

Mutual exclusion (no conflicts — ensures safety/

- Deadlock-freedom (the system as a whole is never

- **Progress** (No subsystem is stuck waiting forever)

Part III Concurrency and Programming Languages

Concurrency and PLs

- Concurrency is about how processes/threads coordinate to achieve a goal.

concrete PL.

change drastically from language to language.

- ... but we don't program in "concurrency", we do it in a
- Concurrency primitives and programming paradigms



errors?

- **Rust:** Resource ownership (compile-time analysis)
- Go: High-level concurrency realized via channels and "goroutines".
- Erlang: High-level concurrency realized with the actor model.

Concurrency and PLs

Can / how do languages help in achieving correctness / ruling out

Concurrency and PLs

Rust:

- No garbage collector (a la C/C++).
- Memory safety (**not** a la C/C++!).
- Memory is managed through a system of **ownership** with a set of rules that the **compiler checks**:
 - Each value in Rust has an owner
 - There can be only one owner at a time
 - When the owner goes out of scope, the value is be dropped.



Rust:

- Ownership can be **borrowed**.

- (Im)mutable references generate (im)mutable borrows.

- Everything else is a **compiler error**.

- At any given time, you can have either one mutable reference or any number of immutable references.



Rust:

- At any given time, you can have either one mutable reference or any number of immutable references.

- Sounds a lot like the discipline to prevent dataraces in shared memory concurrency, doesn't it? More later :)



Go:

- Designed by Google for systems / cloud programming.
- Implementation is garbage-collected.
- Channel-based concurrency built-in.
- (Typed) channels are first-class objects.
- Lightweight threads built-in (goroutines)

Concurrency and PLs

Go:

but not all concurrency woes are solved!

- Pushes developers to synchronize threads using channels.
- Provides various primitives to use and interact with channels.

structuring your code... more later :)

- Threads can read/write to channels concurrently, safely...

- Various concurrency prog. patterns become crucial when



Erlang:

- Aimed at large telecom applications
- Distributed, reliable, soft real-time concurrent systems
- Actor-based concurrency
- Special case of message-passing concurrency



Erlang:

- Actors are isolated computational units

- Communicate by exchanging asynchronous messages

- Messages are queued in a mailbox and processed sequentially.

Concurrency and PLs

- No assumptions on message delivery guarantees.



Erlang:

- Robust failure-handling mechanism.
- Actors can monitor other actors and detect termination.
- "Let it crash" philosophy
- Monitors can restart / kill monitored actors.
- ... more later :)

Concurrency and PLs

Some new "hammers" for specific "nails":

- Go: High-level concurrency realized by "goroutines" that share memory by communicating over channels. High-level / Cloud / Backend programming.

agents and supervision trees. Distributed, reliable, soft real-time concurrent systems.

- **Rust:** Compile-time guarantees of memory safety and race-freedom for shared memory concurrency. Systems / low-level programming.
- Erlang: High-level concurrency realized with the actor model. Actors exchange messages asynchronously. Failure-handling via monitor

Takeaway

- Programming languages should be just like any other tool in your arsenal.
- You should be flexible enough to use the right tool for the job!
- In this course, you will explore three different concurrent programming paradigms and their common idioms.
- By the end, you will have (hopefully) more tools in your toolbox!

Next week:

- Message-passing concurrency module (Go)



- Intro to Go and its message-passing concurrency features

- Some simple programming exercises in Go to get you started.