

# Efficient query processing

Efficient scoring, distributed query processing

Web Search

# Ranking functions

- In general, document scoring functions are of the form

$$\text{score}(q, d) = \text{quality}(d) + \sum_{t \in q} \text{score}(t, d)$$

- The BM25 function, is one of the best performing:

$$\begin{aligned} \text{Score}_{\text{BM25}}(q, d) &= \sum_{t \in q} \log \left( \frac{N}{N_t} \right) \cdot \text{TF}_{\text{BM25}}(t, d), \\ \text{TF}_{\text{BM25}}(t, d) &= \frac{f_{t,d} \cdot (k_1 + 1)}{f_{t,d} + k_1 \cdot ((1 - b) + b \cdot (l_d / l_{\text{avg}}))} \end{aligned}$$

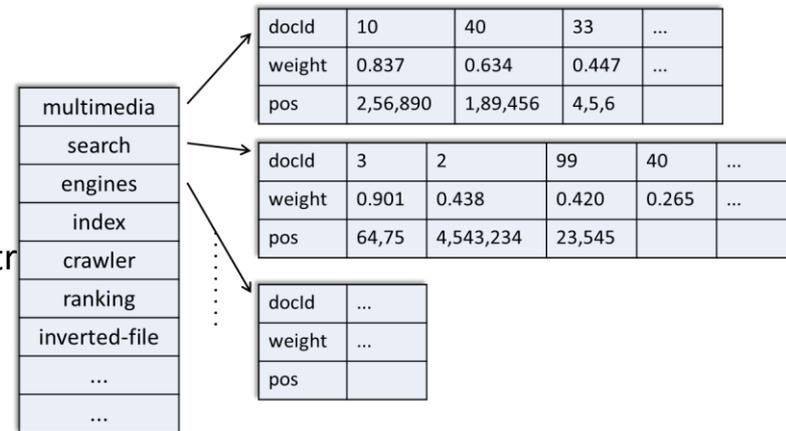
- The term frequency is upper bounded:  $\lim_{f_{t,d} \rightarrow \infty} \text{TF}_{\text{BM25}}(t, d) = k_1 + 1$

# Efficient query processing



## Section 5.1

- Accurate retrieval of top k documents
  - Document at-a-time
  - MaxScore
- Approximate retrieval of top k documents
  - At query time: term-at-a-time
  - At indexing time: term centric and document centric
- Other approaches



# Scoring document-at-a-time

- All documents containing at least one term is scored
- Each document is scored sequentially
  - A naïve method to score all documents is computationally too complex.
- Using a heap to process queries is faster

```
rankBM25_DocumentAtATime ( $\langle t_1, \dots, t_n \rangle, k$ )  $\equiv$   
1    $m \leftarrow 0$  //  $m$  is the total number of matching documents  
2    $d \leftarrow \min_{1 \leq i \leq n} \{\mathbf{nextDoc}(t_i, -\infty)\}$   
3   while  $d < \infty$  do  
4      $results[m].docid \leftarrow d$   
5      $results[m].score \leftarrow \sum_{i=1}^n \log(N/N_{t_i}) \cdot \mathbf{TF}_{\mathbf{BM25}}(t_i, d)$   
6      $m \leftarrow m + 1$   
7      $d \leftarrow \min_{1 \leq i \leq n} \{\mathbf{nextDoc}(t_i, d)\}$   
8   sort  $results[0..(m-1)]$  in decreasing order of  $score$   
9   return  $results[0..(k-1)]$ 
```

**Figure 5.1** Document-at-a-time query processing with BM25.

# Scoring document-at-a-time: Algorithm

**rankBM25\_DocumentAtATime\_WithHeaps**  $((t_1, \dots, t_n), k) \equiv$

Sort in increasing order of score	1 <b>for</b> $i \leftarrow 1$ <b>to</b> $k$ <b>do</b> // create a min-heap for the top $k$ search results 2 $results[i].score \leftarrow 0$
Gets all docs with the query terms	3 <b>for</b> $i \leftarrow 1$ <b>to</b> $n$ <b>do</b> // create a min-heap for the $n$ query terms 4 $terms[i].term \leftarrow t_i$ 5 $terms[i].nextDoc \leftarrow \text{nextDoc}(t_i, -\infty)$ 6 <b>sort</b> $terms$ in increasing order of $nextDoc$ // establish heap property for $terms$
Gets the docs with the lowest ID	7 <b>while</b> $terms[0].nextDoc < \infty$ <b>do</b> 8 $d \leftarrow terms[0].nextDoc$ 9 $score \leftarrow 0$
Process one doc	10 <b>while</b> $terms[0].nextDoc = d$ <b>do</b> 11 $t \leftarrow terms[0].term$ 12 $score \leftarrow score + \log(N/N_t) \cdot \text{TF}_{\text{BM25}}(t, d)$ 13 $terms[0].nextDoc \leftarrow \text{nextDoc}(t, d)$ 14 <b>reheap</b> ( $terms$ ) // restore heap property for $terms$
Replace the worst doc	15 <b>if</b> $score > results[0].score$ <b>then</b> 16 $results[0].docid \leftarrow d$ 17 $results[0].score \leftarrow score$ 18 <b>reheap</b> ( $results$ ) // restore heap property for $results$
Sort in decreasing order of score	19 <b>remove from</b> $results$ all items with $score = 0$ 20 <b>sort</b> $results$ in decreasing order of $score$ 21 <b>return</b> $results$

**Figure 5.3** Document-at-a-time query processing with BM25, using binary heaps for managing the set of terms and managing the set of top- $k$  documents.

# MaxScore

- We know that each term frequency is bounded by

$$\lim_{f_{t,d} \rightarrow \infty} \text{TF}_{\text{BM25}}(t, d) = k_1 + 1$$

- We call this score the MaxScore of a term
- If the score of the k'th document exceeds the MaxScore of a term X,
  - We can ignore documents containing only term X
  - When considering term Y, we still need to check the term X contribution
  - If the score of the k'th document exceeds the MaxScore of terms X and Y,
    - We can ignore documents containing terms Y

# Scoring document-at-a-time: comparison

- Comparison between reheap with & w/out MaxScore

**Table 5.1** Total time per query and CPU time per query, with and without MAXSCORE. Data set: 10,000 queries from TREC TB 2006, evaluated against a frequency index for GOV2.

	Without MaxScore			With MaxScore		
	Wall Time	CPU	Docs Scored	Wall Time	CPU	Docs Scored
OR, k=10	400 ms	304 ms	$4.4 \cdot 10^6$	188 ms	93 ms	$2.8 \cdot 10^5$
OR, k=100	402 ms	306 ms	$4.4 \cdot 10^6$	206 ms	110 ms	$3.9 \cdot 10^5$
OR, k=1000	426 ms	329 ms	$4.4 \cdot 10^6$	249 ms	152 ms	$6.2 \cdot 10^5$
AND, k=10	160 ms	62 ms	$2.8 \cdot 10^4$	n/a	n/a	n/a

**Both methods are exact!**

# Approximating the $K$ largest scores

- Typically we want to retrieve the top  $K$  docs
  - not to totally order all docs in the collection
- Can we approximate the  $K$  highest scoring documents?
- Let  $J$  = number of docs with nonzero scores
  - We seek the  $K$  best of these  $J$

# Scoring term-at-time

- The index is organized by postings-lists
  - Processing a query a document-at-a-time requires several disk seeks
  - Processing a query a term-at-a-time minimizes disk seeks
- In this method, a query is processed a term-at-a-time and an accumulator stores the score of each term in the query.
- When all terms are processed, the accumulator contains the scores of the documents.
- Do we need to have an accumulator the size of the collection or the largest posting list?

# Scoring term-at-time

- A query is processed a term-at-a-time and an accumulator stores the score of each term in the query.
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```
rankBM25_TermAtATime ( $(t_1, \dots, t_n), k$ )  $\equiv$ 
1  sort  $\langle t_1, \dots, t_n \rangle$  in increasing order of  $N_{t_i}$ 
2   $acc \leftarrow \{\}, acc' \leftarrow \{\}$  // initialize two empty accumulator sets
3   $acc[0].docid \leftarrow \infty$  // end-of-list marker
4  for  $i \leftarrow 1$  to  $n$  do
5       $inPos \leftarrow 0$  // current position in  $acc$ 
6       $outPos \leftarrow 0$  // current position in  $acc'$ 
7      for each document  $d$  in  $t_i$ 's postings list do
8          while  $acc[inPos] < d$  do // copy accumulators from  $acc$  to  $acc'$ 
9               $acc'[outPos++] \leftarrow acc[inPos++]$ 
10              $acc'[outPos].docid \leftarrow d$ 
11              $acc'[outPos].score \leftarrow \log(N/N_{t_i}) \cdot TF_{BM25}(t_i, d)$ 
12             if  $acc[inPos].docid = d$  then // term and accumulator coincide
13                  $acc'[outPos].score \leftarrow acc'[outPos].score + acc[inPos].score$ 
14                  $d \leftarrow \text{nextDoc}(t_i, acc'[outPos])$ 
15                  $outPos \leftarrow outPos + 1$ 
16             while  $acc[inPos] < \infty$  do // copy remaining accumulators from  $acc$  to  $acc'$ 
17                  $acc'[outPos++] \leftarrow acc[inPos++]$ 
18                  $acc'[outPos].docid \leftarrow \infty$  // end-of-list marker
19             swap  $acc$  and  $acc'$ 
20     return the top  $k$  items of  $acc$  // use a heap to select the top  $k$ 
```

**Figure 5.4** Term-at-a-time query processing with BM25. Document scores are stored in accumulators. The accumulator array is traversed co-sequentially with the current term's postings list.

# Limited accumulator

lec 7.5

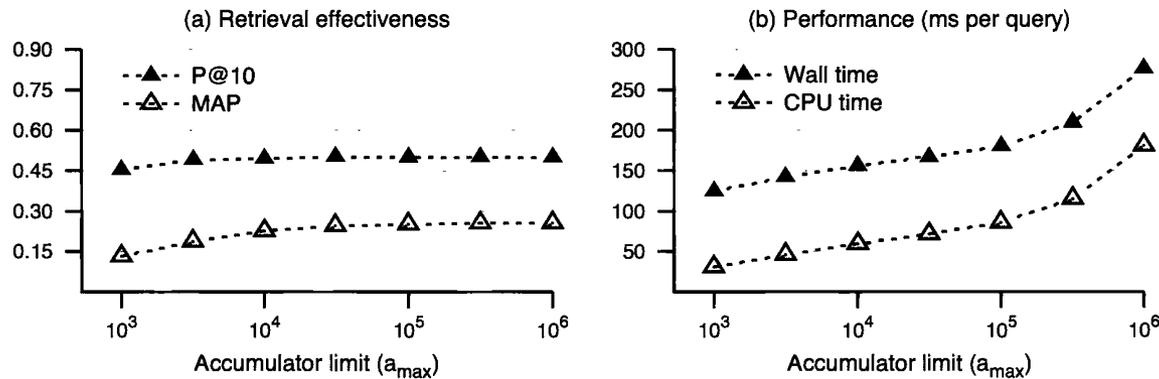
- The accumulator may not fit in memory, so, we ought to limit the accumulator's length
- When traversing  $t$ 's postings
  - Add the posting only if it is below a  $v_{TF}$  threshold
- For each document in the postings list
  - accumulate the term score or use new positions in accumulator for that doc

# High-idf query terms only

Section 5

- When considering the postings of query terms
- Look at them in order of decreasing idf
  - High idf terms likely to contribute most to score
- For a query such as “catcher in the rye”
  - Only accumulate scores from “catcher” and “rye”

# Scoring term-at-a-time

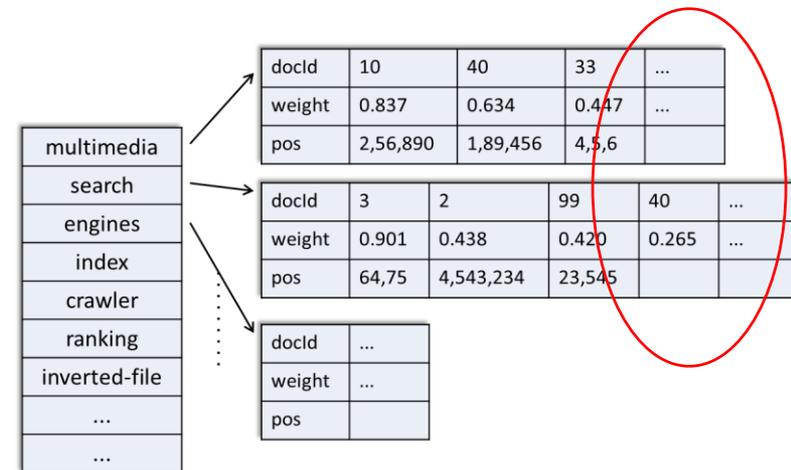


**Figure 5.6** Retrieval effectiveness and query processing performance for term-at-a-time query evaluation with accumulator pruning. Data set: 10,000 queries from TREC TB 2006, run against a frequency index for GOV2.

- Baseline:
  - Top 10 MaxScore 188ms, 93 ms,  $2.8 \times 10^5$  docs
  - Top 1000 MaxScore 242ms, 152 ms,  $6.2 \times 10^5$  docs

# Index pruning

- The accumulator technique ignores several query term's postings
  - This is done in query time.
- How can we prune postings that we know in advance that they will be almost noise?
- The goal is to keep only the most informative postings in the index.

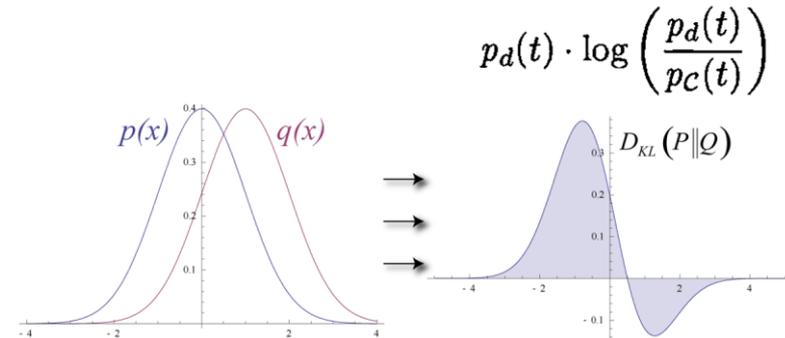


# Term-centric index pruning

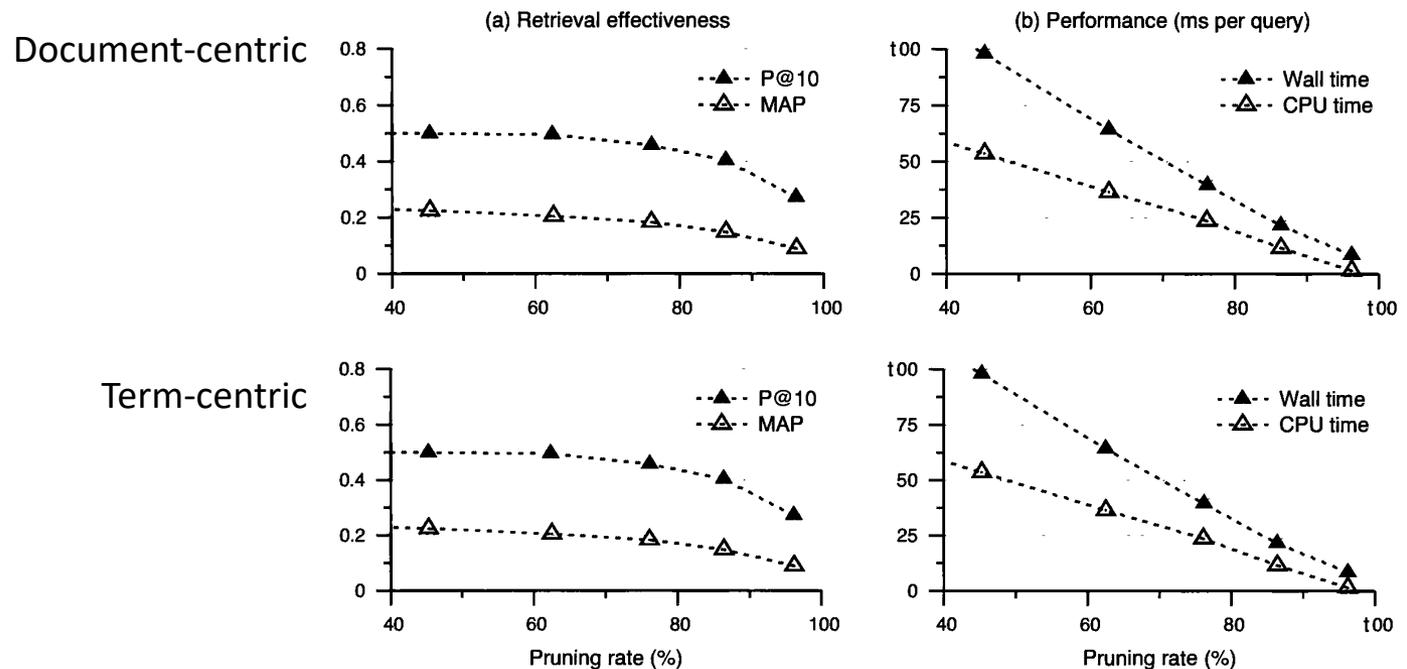
- **Examining only term postings**, we can decide if a given term is relevant in general (IDF) or relevant for the document (TF).
- If a term appears less than  $K$  times in documents, store all of  $t$ 's postings in the index
- If the term  $t$  appears in more than  $K$  documents
  - Compute the term score of the  $K$ 'th document
  - Consider only the postings with scores  $> score(d_k, t) \cdot \epsilon$  where  $0.0 < \epsilon < 1.0$

# Document-centric index pruning

- **Examining each document's terms distribution** we can predict which terms are the most representative of that document.
  - The terms is added to the index only if it is considered representative of the document
- Compute the Kullback-Leibler divergence between the terms distribution in the document and in the collection.
- For each document, select the top  $\lambda \cdot n$  terms to be added to the index

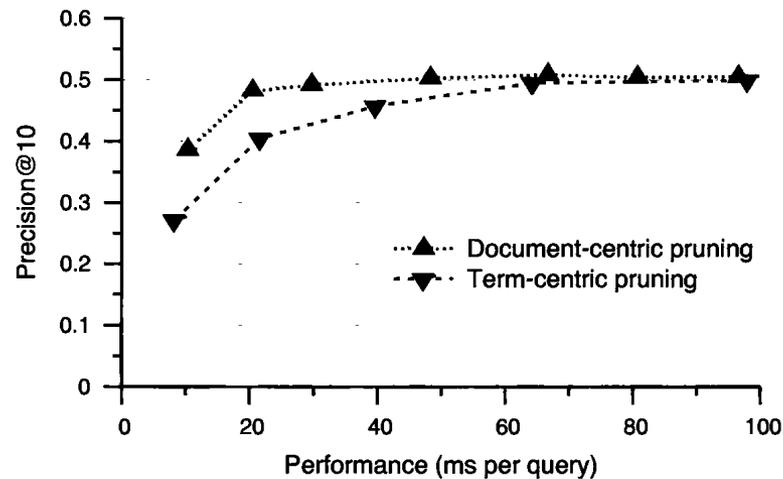


# Head-to-head comparison



**Figure 5.8** Term-centric index pruning with  $K = 1,000$  and  $\epsilon$  between 0.5 and 1. Data set for efficiency evaluation: 10,000 queries from TREC TB 2006. Data set for effectiveness evaluation: TREC topics 701-800.

# Head-to-head comparison



- Baseline:

- Top 10

MaxScore 188ms, 93 ms,  $2.8 \times 10^5$  docs

- Top 1000

MaxScore 242ms, 152 ms,  $6.2 \times 10^5$  docs

# Other approaches



- Static scores
- Cluster pruning
- Number of query terms
- Impact ordering
  - Champion lists
  - Tiered indexes

# Static quality scores

- We want top-ranking documents to be both relevant and authoritative
- Relevance is being modeled by cosine scores
- Authority is typically a query-independent property of a document
- Examples of authority signals
  - Wikipedia among websites
  - Articles in certain newspapers
  - A paper with many citations
  - Many diggs, Y!buzzes or del.icio.us marks
  - (Pagerank)

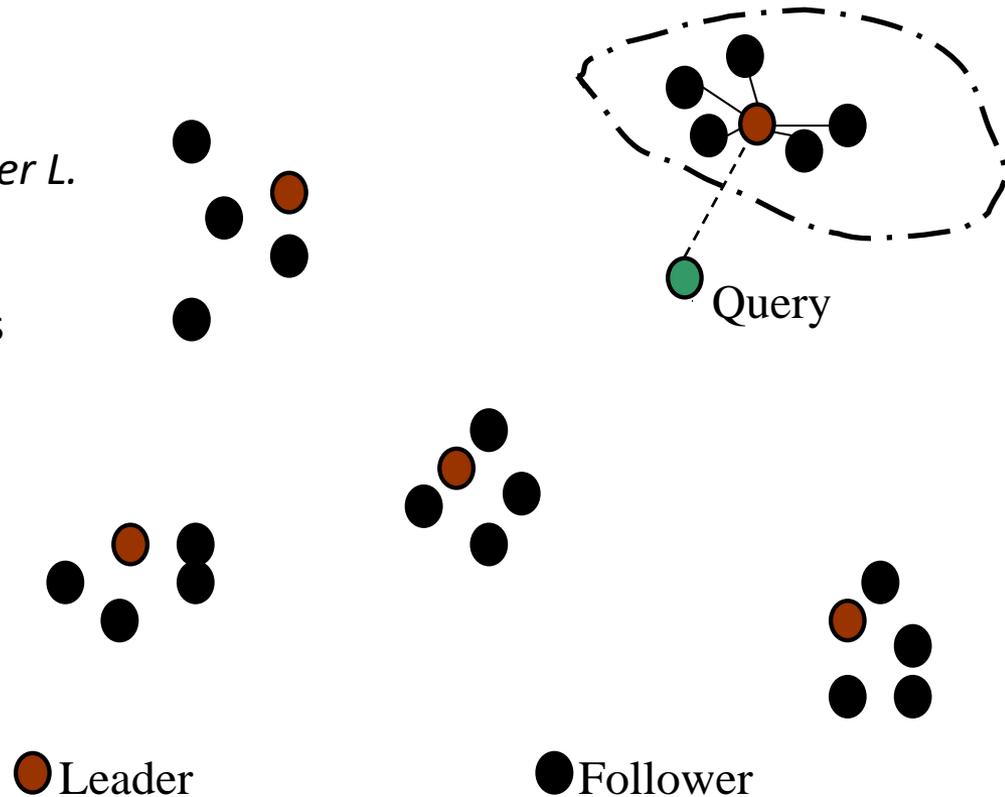


# Cluster pruning: preprocessing

- Pick  $\sqrt{N}$  docs at random: call these *leaders*
- For every other doc, pre-compute nearest leader
  - Docs attached to a leader: its *followers*;
  - Likely: each leader has  $\sim \sqrt{N}$  followers.

# Cluster pruning: query processing

- Given query  $Q$ , find its nearest *leader*  $L$ .
- Seek  $K$  nearest docs from among  $L$ 's followers.



# Champion lists

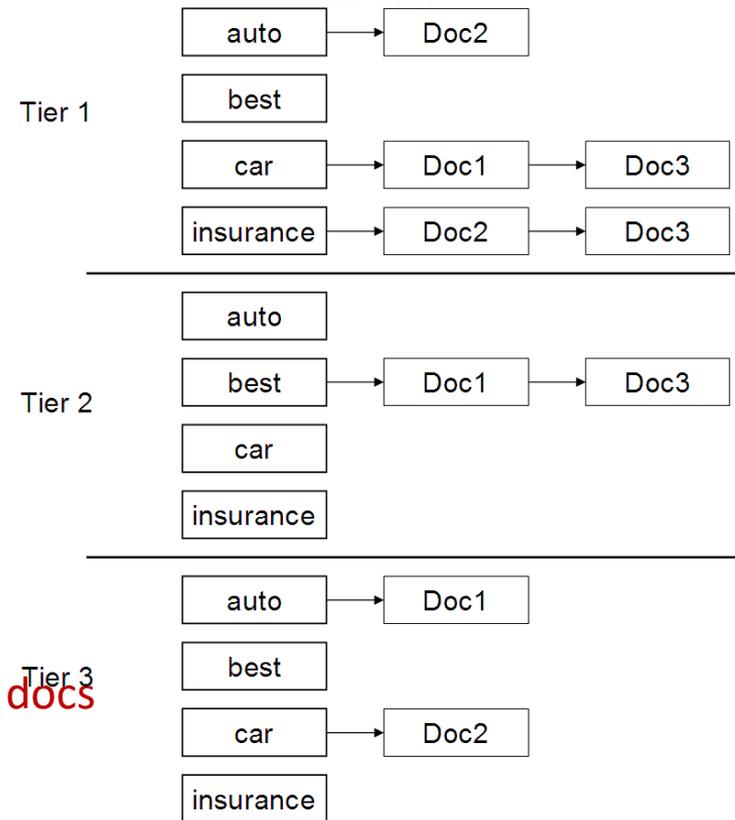
- Precompute for each dictionary term  $t$ , the  $r$  docs of highest weight in  $t$ 's postings
  - Call this the champion list for  $t$
  - (aka fancy list or top docs for  $t$ )
- Note that  $r$  has to be chosen at index build time
  - Thus, it's possible that  $r < K$
- At query time, only compute scores for docs in the champion list of some query term
  - Pick the  $K$  top-scoring docs from amongst these

# Docs containing many query terms

- Any doc with at least one query term is a candidate for the top  $K$  output list
- For multi-term queries, only compute scores for docs containing several of the query terms
  - Say, at least 3 out of 4
  - Imposes a “soft conjunction” on queries seen on web search engines (early Google)
- Easy to implement in postings traversal

# Tiered indexes

- Break postings up into a hierarchy of lists
  - Most important
  - ...
  - Least important
- Inverted index thus broken up into tiers of decreasing importance
- **At query time use top tier unless it fails to yield  $K$  docs**
  - If so drop to lower tiers
  - Common practice in web search engines



# Scalability: Index partitioning



**Document partitioning**

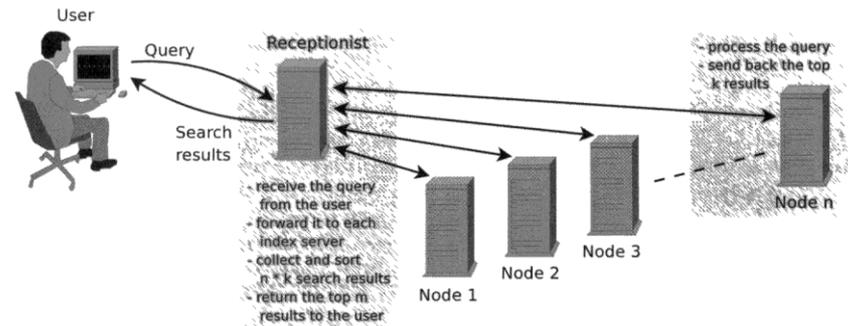
		Documents								
		D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	D <sub>5</sub>	D <sub>6</sub>	D <sub>7</sub>	D <sub>8</sub>	D <sub>9</sub>
Terms	T <sub>1</sub>	X		X	X		X			X
	T <sub>2</sub>		X			X				
	T <sub>3</sub>		X	X					X	
	T <sub>4</sub>				X			X		
	T <sub>5</sub>	X					X			X
	T <sub>6</sub>	X						X	X	
	T <sub>7</sub>		X		X		X			
	T <sub>8</sub>			X						X
		Node 1			Node 2			Node 3		

**Term partitioning**

		Documents								
		D <sub>1</sub>	D <sub>2</sub>	D <sub>3</sub>	D <sub>4</sub>	D <sub>5</sub>	D <sub>6</sub>	D <sub>7</sub>	D <sub>8</sub>	D <sub>9</sub>
Terms	T <sub>1</sub>	X		X	X		X			X
	T <sub>2</sub>		X			X				
	T <sub>3</sub>		X	X					X	
	T <sub>4</sub>				X			X		
	T <sub>5</sub>	X					X			X
	T <sub>6</sub>	X						X	X	
	T <sub>7</sub>		X		X		X			
	T <sub>8</sub>			X						X
		Node 1			Node 2			Node 3		

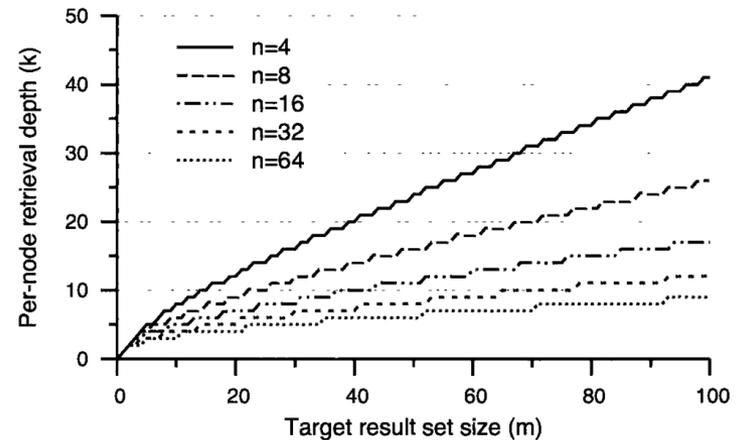
# Doc-partitioning indexes

- Each index server is responsible for a random sub-set of the documents
- All  $n$  nodes return  $k$  results to produce the final list with  $m$  results
- Requires a background process to keep the *idf* (and other general statistics) synchronized across index servers



# How many documents to return per index-server?

- The choice of  $k$  has impact on:
  - The network load to transfer partial search results from the index-servers to the server doing the rank fusion;
  - The precision of the final rank.



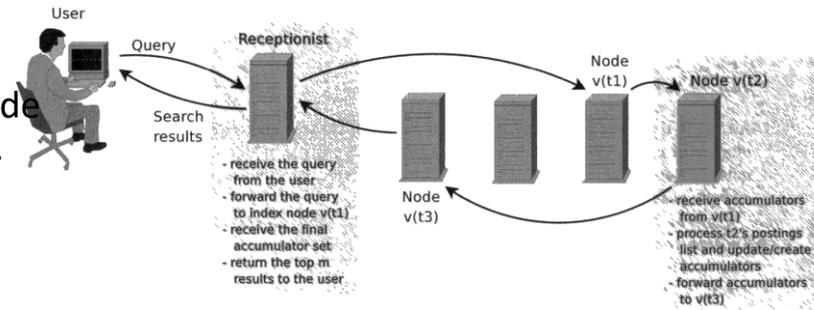
**Figure 14.3** Choosing the minimum retrieval depth  $k$  that returns the top  $m$  results with probability  $p(n, m, k) > 99.9\%$ , where  $n$  is the number of nodes in the document-partitioned index.

# Term-partitioning indexes

- Each index server receives a sub-set of the dictionary's terms

- A query is sent simultaneously to the term's corresponding nodes.

- Each node passes its accumulator to the next node or to the central node to compute the final rank.



- Disadvantages:

- This requires that each node loads the full posting list for each term.
  - Uneven load balance due to query drifts.
  - Unable to do support efficient document-at-a-time scoring.

# Planet-scale load-balancing

- When a systems receives requests from the entire planet at every second...
- The best way to load-balance the queries is to use DNS to distributed queries across data-centers.
- Each query is assigned a different IP according to the data-center load and to the user's geographic location.

## Serving a Google query

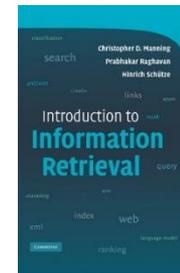
When a user enters a query to Google (such as `www.google.com/search?q=ieee+society`), the user's browser first performs a domain name system (DNS) lookup to map `www.google.com` to a particular IP address. To provide sufficient capacity to handle query traffic, our service consists of multiple clusters distributed worldwide.

Each cluster has around a few thousand machines, and the geographically distributed setup protects us against catastrophic data center failures (like those arising from earthquakes and large-scale power failures). A DNS-based load-balancing system selects a cluster by accounting for the user's geographic proximity to each physical cluster. The load-balancing system minimizes round-trip time for the user's request, while also considering the available capacity at the various clusters.

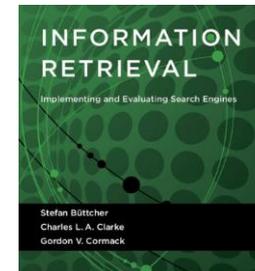
Barroso, Luiz André, Jeffrey Dean, and Urs Hölzle. "Web search for a planet: The Google cluster architecture." *IEEE Micro* (2003)

# Summary

- Relevance feedback
  - Pseudo-relevance feedback
- Query expansion
  - Dictionary based
  - Statistical analysis of words co-occurrences
- Efficient scoring
  - Per-term and per-doc pruning
- Distributed query processing
  - Per-term and per-doc pruning



Chapter 7 and 9



Section 5.1

Section 14.1